



Ongoing evolution of Linux x86 machine check handling

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What's a good error?

- User has to see it, of course
 - That can be surprisingly difficult
 - Also psychological barriers (“users don't read errors”)
- High level classification
 - Software error versus hardware error
 - Don't want a hardware error reported as a kernel bug
 - Still have low level details for experts (ideally separated)
- Identify affected component
 - Do not require low level knowledge to process
 - Works out of the box
- How serious is the error?
 - Do not upset the user unnecessarily

Error sources

- Machine checks from the CPU
- NMI
- PCI-Express Advanced Error Handling (AER)
- Chipset
- ACPI4 (APEI)
- Drivers
 - SATA errors
 - Ethernet
 - ...
- Presentation only covers machine check errors from the CPU
 - This includes memory on modern systems

What is a machine check?

- Machine check is a hardware error reported by the CPU
 - Not a software problem!
- Hardware corrects most problems, but sometimes it can fail
 - Memory, Interconnect, Cache, Internal errors
- Uncorrected errors raise exceptions (“MCEs”)
 - Better to stop than to continue with corrupted data
 - Otherwise the corrupted data could hit disk or give wrong results
 - If you aren't sure where it is, stop the machine ...
- Corrected errors are reported in the background
 - Either using a poll handler or with interrupts (“CMCI”)
 - Didn't cause corruption (yet)

Why are machine checks important?

- They report memory errors on modern systems
 - Memory error rate scaling roughly with memory sizes
 - Memory sizes are increasing quickly
 - More cores need more memory
 - Virtualization needs a lot of memory
 - This also means more memory errors
 - So good error handling is important
- On large clusters, errors are common
 - What's uncommon on a single system
 - ... becomes common when you have a hundred of them
 - ... and even very rare events become common on thousands of systems
- In general, good error diagnosis is useful
 - If you ever searched manually for a bad DIMM ...
 - Saves time and hassle

MCE errors in practice today

- Error flows for uncorrected and corrected errors
- Assuming 64 bit or 2.6.31+ with CONFIG_X86_NEW_MCE
- Mcelog has to be installed
- Some of these flows are still somewhat clumsy
 - The future will be brighter, hopefully

Classic unrecoverable MCE error today

- System detects uncorrected error
 - Requires ECC for memory
 - Machine check exception happens
- Machine check handler collects error and prints it out
 - CPU x: Machine check exception ...
- System panics on unrecoverable error
 - On auto reboot (panic=30) can be logged after reboot on many systems
 - Available then decoded in /var/log/mcelog some time after boot
 - Trap: doesn't work with double reboot or with power switch
 - You will see the panic on the console (not in X) or on serial/netconsole
 - If you don't have auto reboot a logging console is very useful
 - Console output can be run through mcelog --ascii to decode
- Analyze error, based on decoded output
 - For example, map to DIMM (“Memory DIMM ID of error: ...”)
 - If common, take corrective action

Corrected error flow today

- A data bit flips
 - Hardware detects error, using checksum, and corrects it and reports event
- CMCi interrupt happens or poll timer picks it up
 - Kernel logs it to internal buffer accessible through /dev/mcelog
 - Note this buffer can overflow
- Mcelog picks it up
 - Either as cronjob or in daemon mode or as trigger
 - In cronjob mode up to 10 minutes delay, worst case with polling
 - Mcelog decodes
 - Logged in /var/log/mcelog or syslog
- Analysis of the log entry
 - Identify component
 - Take corrective action if common

Low level MCE handler improvements

- General overhaul after comprehensive audit
 - A lot of small improvements, too many to list
- Monarch support: synchronization over all CPUs
 - Collect errors from all CPUs
 - Synchronize all CPUs
 - Process the most serious error first to avoid data corruption
 - When the hardware didn't contain the error shut down
- Bank sharing
 - Handle shared machine check banks correctly
- Corrected Machine Check Interrupt Support (“CMCI”)
- Injector support for testing and a comprehensive test suite

MCA recovery

- New CPU feature in upcoming Nehalem-EX CPU
- Recovery from some memory uncorrected errors
 - For example, Patrol scrub memory error in the background
 - Required a lot of changes in the MCE handler to do reliable
 - When recovering it's much more important to handle all corner cases
- OS finds out what the corrupted page does
 - And attempts to get rid of it
- Machine check architecture has new status bits for recovery
 - Signalled, Action-Required
 - Different types of errors: Action-Optional, Action-Required, UCNA, Corrected

HWPoison handling in the VM

- VM finds out who owns a page and stops using it
 - Pages with copy on disk can be just dropped
 - Or application is killed, if data has no copy
 - IO error for dirty file cache pages
 - Free pages will be ignored on allocation
- Difficult because error can come in at any time
 - Can disrupt normal page lifecycle
 - Error code has to be very careful
 - Also testing is difficult
- Put page on bad page list and never reuse it again
- Page table entry of any mappings is poisoned
- Early kill versus late kill mode

Virtualization

- Virtualization requires a lot of memory
- Many eggs in a single basket (“servers on a single machine”)
- Error handling and error containment is important
 - If there's a problem only kill single guest, not all
- KVM guests act like processes
 - So uses the per process infrastructure in standard hwpoison
- Uncorrected recoverable errors can be forwarded to a guest
 - Guest can inject the error as machine check (or KVM kills)
 - Guest with MCA recovery support can recover (or panic)
- Similar code developed for Xen
 - No forwarding for Xen currently



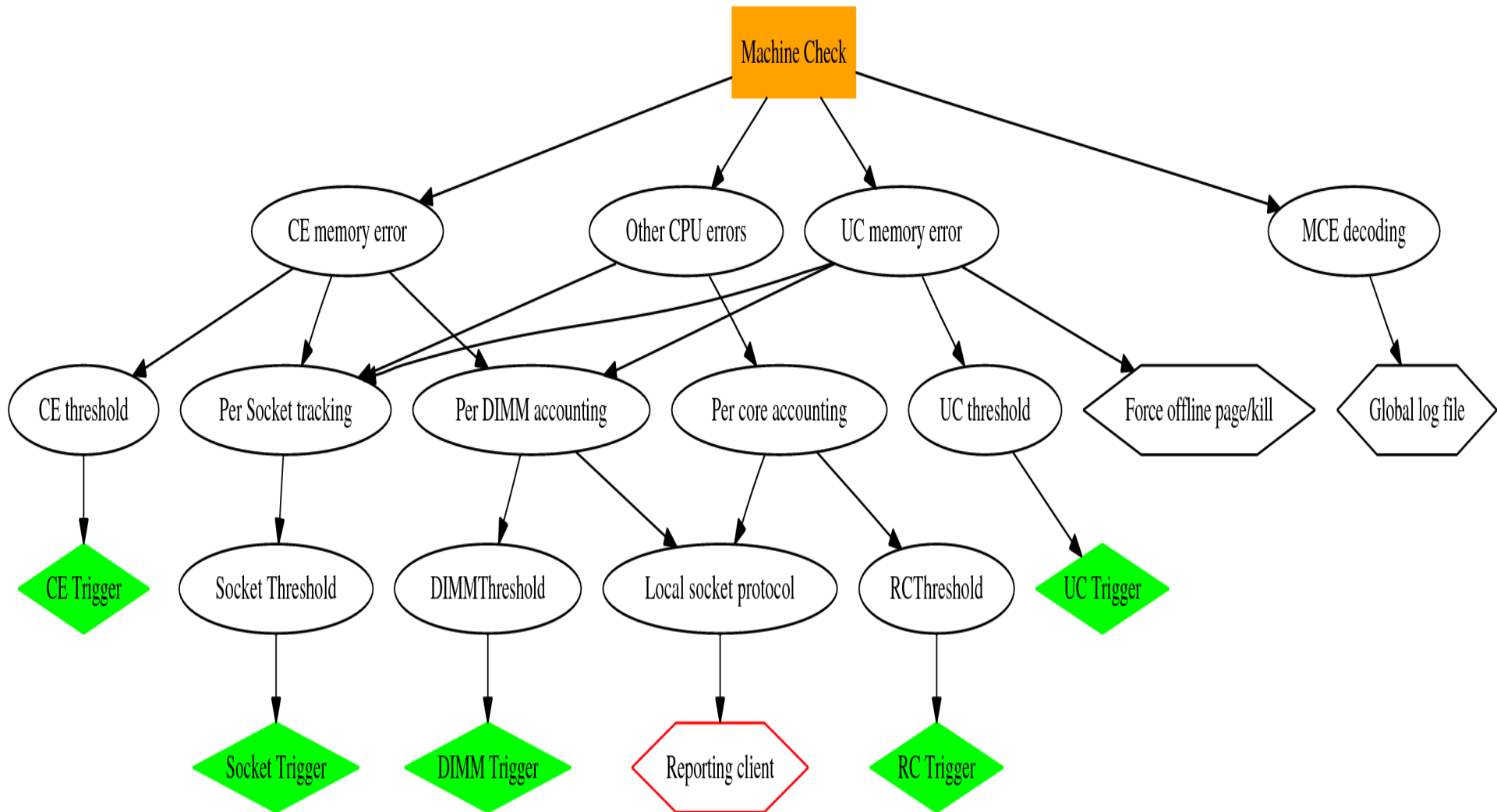
Memory error application interface

- Applications can catch memory errors, which are signals
 - Was needed for KVM, but can be used by others too
- Applications have often cache, which they can drop
- Expect only a few specialized applications to use it
- SIGBUS with an address can be caught
 - BUS_MCEERR_AO Action-Optional error in process
 - Get rid of page specified by si_addr, si_addr_lsb
 - Take out of free list or similar
 - BUS_MCEERR_AR Action-Required error in current execution context
 - Need to abort right now (siglongjmp etc.)
- Prctl to set early kill vs late kill for the process
 - Late kill is typically better for error aware applications

Mcelog

- User space backend that decodes and processes MCE errors
 - Also identifies components with some firmware help
- Traditionally on 64 bit x86, now on 32 bit too
- Traditionally cronjob every 5 minutes, future daemon
 - Daemon mode allows to keep state about errors in memory with query interface and triggers
 - Some old attempts using an on-disk database proved difficult
- Can run shell scripts on specific events (“triggers”)
 - Notify administrator, offline component, ...
- Long term goal: high level errors in syslog
 - Some steps into this direction

Mcelog error flow



Error accounting

- Possible in mcelog daemon mode
- Often most interesting is which component the error affects
 - DIMM, memory, PCI card, etc.
 - To see trends and replace the right component quickly if needed
- Individual errors are often not that interesting
 - Errors often come in bursts and individual errors in a burst are not interesting
 - Large clusters can generate a lot of data, which is difficult to process
- Mcelog moving towards accounting errors per component
 - Only reports “n errors in last x hours on component k”
 - Triggers when thresholds are exceeded
 - Discovers component names with firmware help
 - Can disable individual error logging for less data

Open issues

- Crashdump handling
- More testing is always useful
 - Stress test suite under way
 - Any contributions welcome
- Better error reporting in general
 - More high level errors better presented
- More error sources in mcelog
- Intelligent error handling in mcelog
 - If you have ideas, feel free to contact me

Resources

- <http://halobates.de/mce.pdf>
 - Old paper about Linux machine checks
- Intel Software Developer's Manual: 3A/B System Programming Guide
 - Description of the x86 Machine check architecture
- [git://git.kernel.org/pub/scm/linux/kernel/git/ak/linux-mce-2.6](https://git.kernel.org/pub/scm/linux/kernel/git/ak/linux-mce-2.6)
 - MCE development tree (hwpoison, mce4)
- [git://git.kernel.org/pub/scm/utils/cpu/mcelog.git](https://git.kernel.org/pub/scm/utils/cpu/mcelog.git)
 - Mcelog development repository
- <ftp://ftp.kernel.org/pub/linux/utils/cpu/mce/>
 - Mcelog releases
- [git://git.kernel.org/pub/scm/utils/cpu/mce-test.git](https://git.kernel.org/pub/scm/utils/cpu/mce-test.git)
 - MCE test suite (or now part of LTP to (<http://ltp.sourceforge.net>))
- [git://git.kernel.org/pub/scm/utils/cpu/mce-inject.git](https://git.kernel.org/pub/scm/utils/cpu/mce-inject.git)
 - MCE injector



Backup



Mcelog versus EDAC

- EDAC old style driver for chipset memory controllers
 - Exposes a lot of low level details
 - New model: memory controller in CPU
 - Memory errors integrated with machine checks
 - Handled by standard MCE subsystem
- EDAC needs driver for each platform
 - And often accesses “non stable” registers that could change even with steppings
 - Mcelog uses standardized interfaces
- No integration with software
 - Requires special configuration for each board to identify components
- Cannot do a lot of things that user space (mcelog) can do

Testing

- Testing machine checks is difficult
 - Normal operation doesn't have enough errors
 - So standard Linux community testing model doesn't work very
 - Needs special injection support and test suites
- Various injectors on software level (without hardware support)
 - Low level machine check injector
 - Page error injector for process and for arbitrary page
- New injectors, test suite mce-test for testing MCEs
 - Testing low level handler with mce-inject
 - Testing hwpoison VM code in process context
 - Bring pages into specific states and test to see if they can be poisoned
- Ongoing work to get the best test coverage

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